



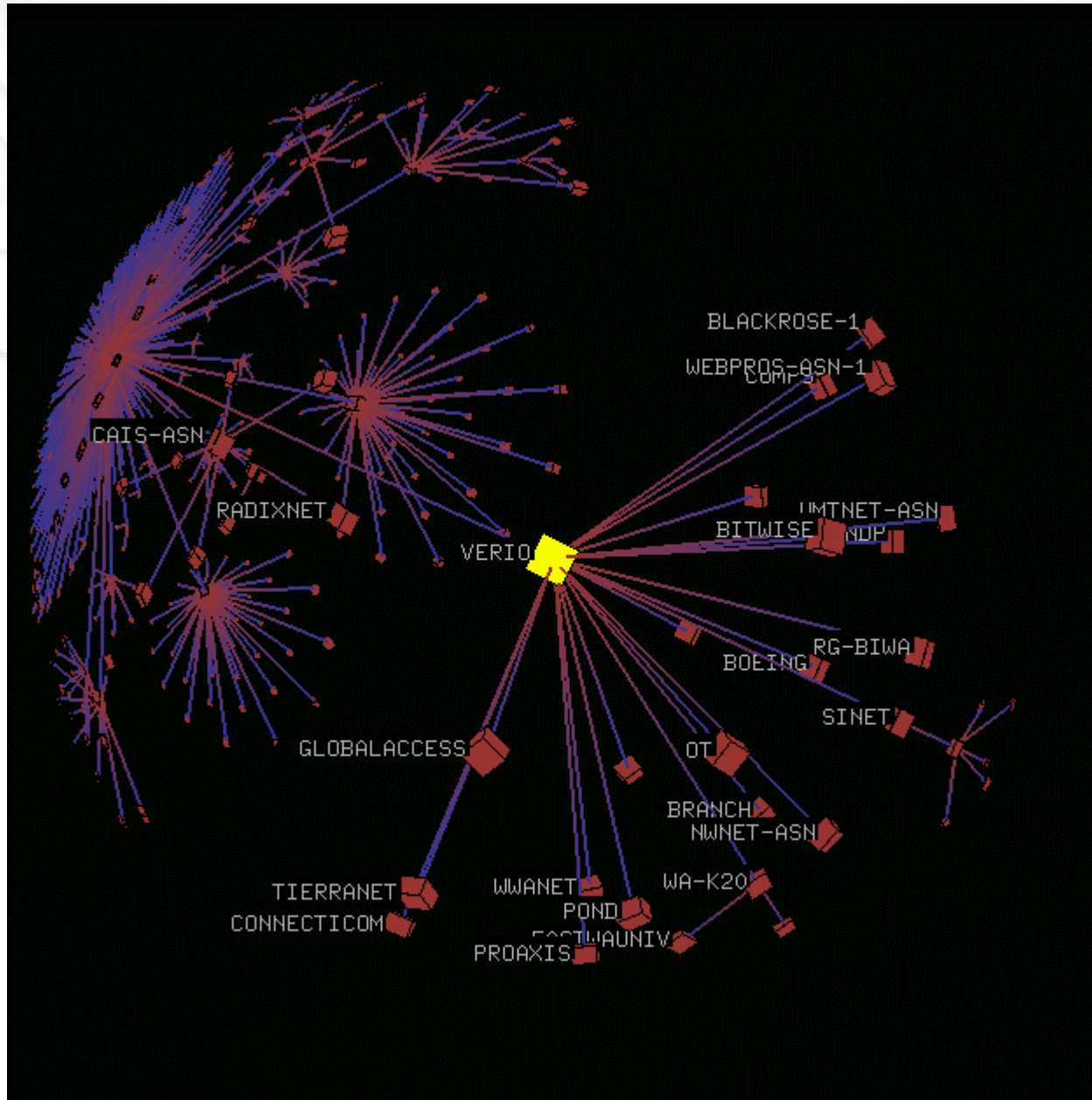
Fractals in Networking: Modeling and Inference

Rolf Riedi

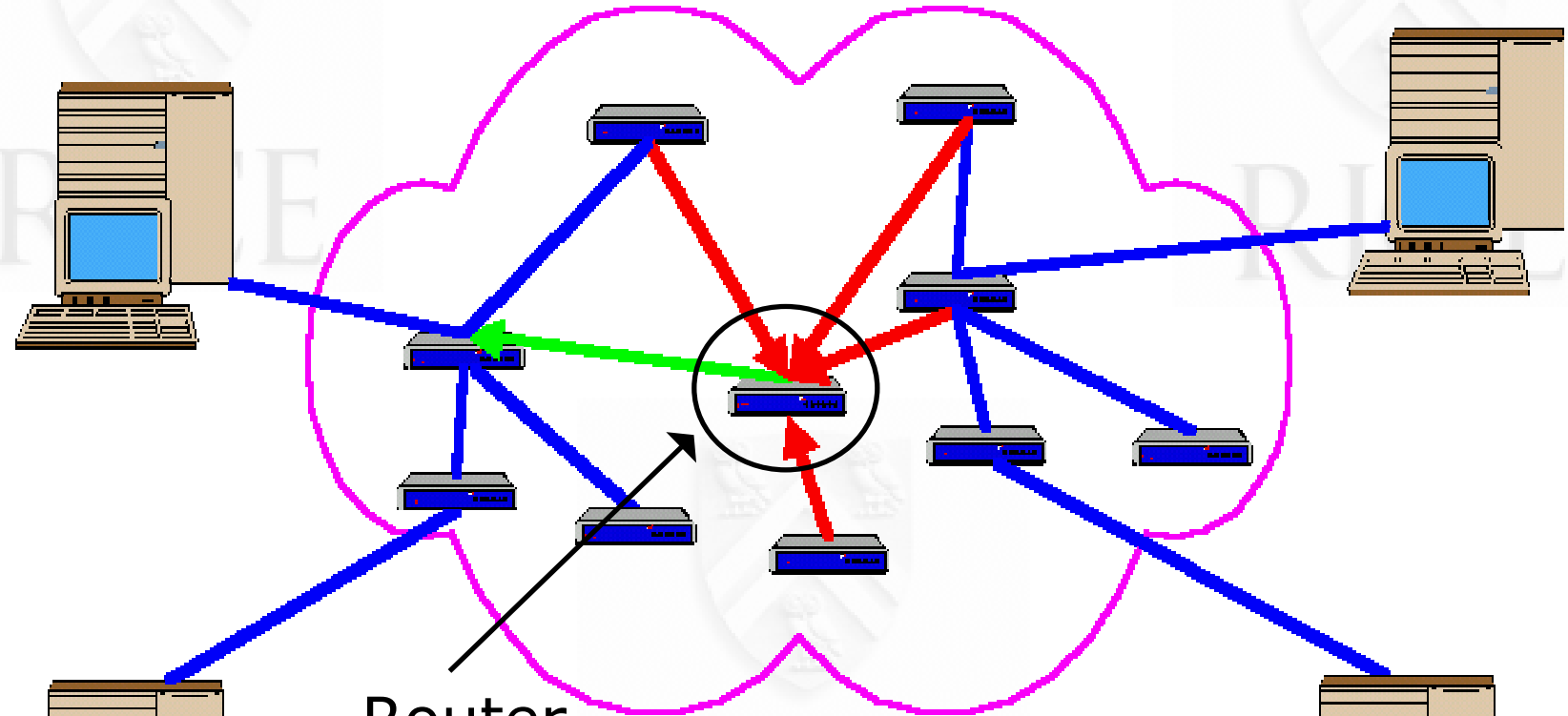
R. Baraniuk

A. Keshavarz-Haddad, V. Ribeiro, S. Sarvotham
spin.rice.edu

Fractal2004, Vancouver, April 2004



Internet is packet switched

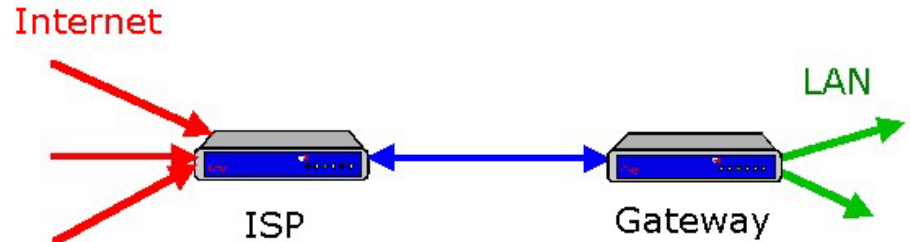


Router

- Data packets are switched to the outgoing port
- there they are **queued**
- ...or **dropped**.

Measured Data

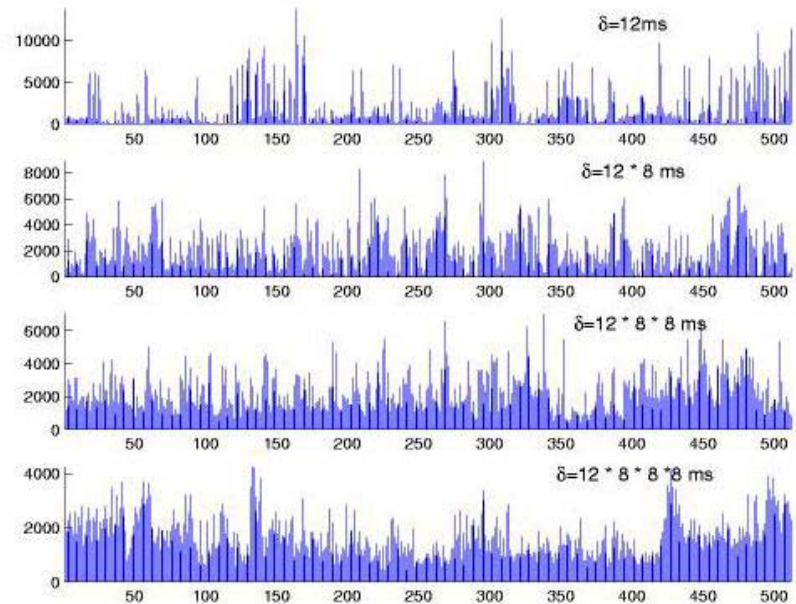
- Time series (A_k, Z_k) collected at gateway of LAN
 - k = number of data packet
 - A_k = arrival time of packet
 - B_k = size of packet



- Working data:

Bytes arriving in intervals of m time units

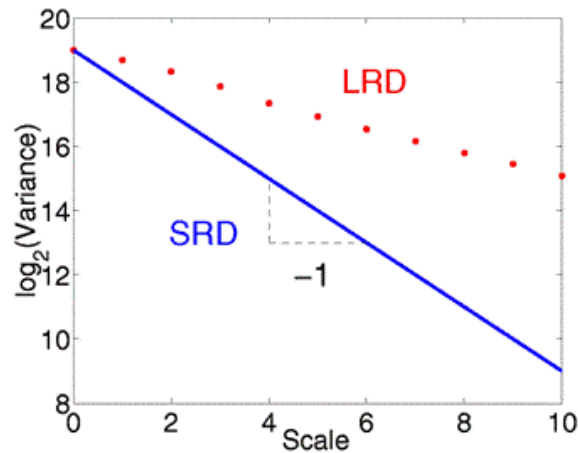
$$X_n^{(m)} = \sum_{mn+1 < A_k < m(n+1)} B_k$$



- Observation/Motivation:

High variability of $X_n^{(m)}$ for large m degrading performance (loss, delay)

Network traffic is LRD (Long Range Dependent)

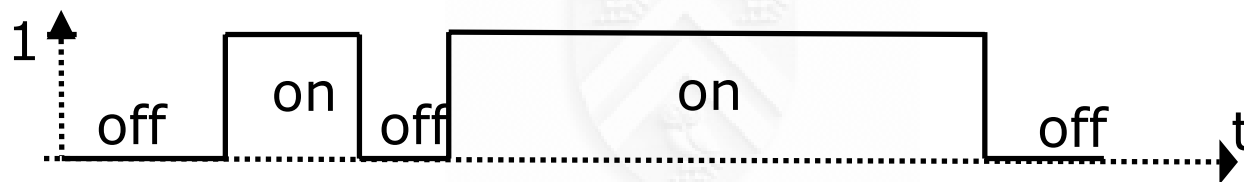


Bellcore '89

Network traffic is LRD

- Failure of Poisson modeling (Paxson Floyd 1995)
- Discovery of LRD -- Mathematical model

Source i :
$$X_i(t) = \begin{cases} 1 & \text{if time } t \text{ is an } \textit{on} \text{ interval,} \\ 0 & \text{if time } t \text{ is in } \textit{off} \text{ interval.} \end{cases}$$



Heavy tailed on- and off-durations, e.g.,

$$1 - F_{on}(x) \sim \ell_{on} x^{-\alpha_{on}}, \quad 1 < \alpha_{on} < 2$$

(Willinger Taqqu Wilson Leland '93, Renewal reward: Mandelbrot 60's)

Limiting behavior

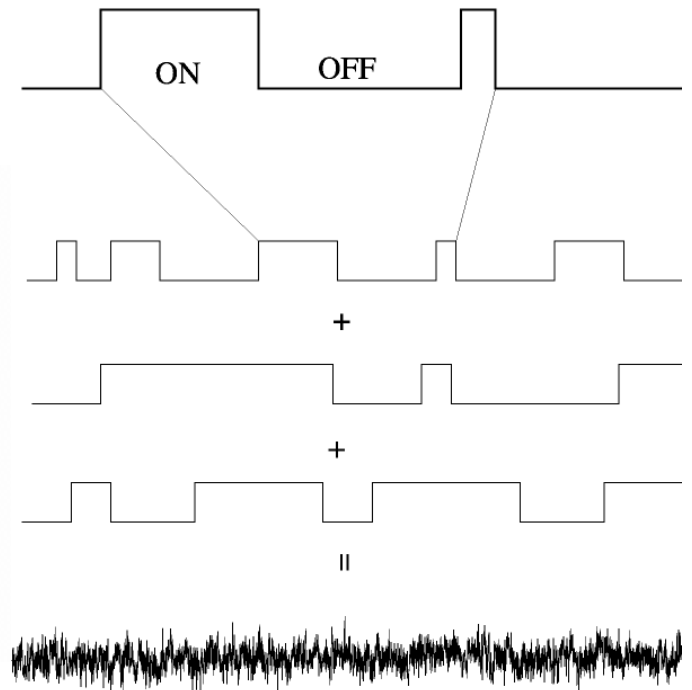
Multiplexed Traffic:
$$W^{(m)}(t) = \sum_{i=1}^m X_i(t)$$

$$\frac{W^{(m)}(t) - \mathbb{E}W^{(m)}}{m^{1/2}} \xrightarrow{fdd} G(t)$$

$$\frac{1}{TH} \int_0^{Tt} G(u) du \xrightarrow{fdd} \sigma B_H(t)$$

$$H = \frac{3 - \min(\alpha_{on}, \alpha_{off})}{2}$$

(Lamperti 1962)
(Taqqu Levy 1986)



Physical model

- On-off model: physically appealing

- Verification

- Heavy tailed file sizes on web (Crovella '96)
- Client behavior

- Implications

- Using self-similarity: (Norros '94)

To reduce overflow increase bandwidth, not buffer

- Using large deviations: (Duffield '95)

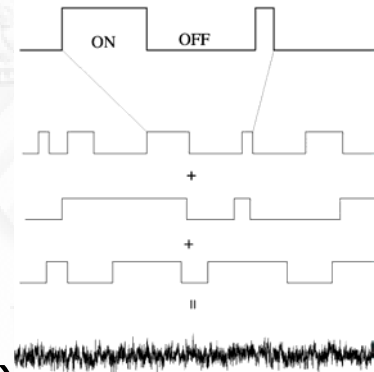
Queue-tail is Weibull $P[Q > b] \simeq \exp(-b^{2H})$

- Predictive control

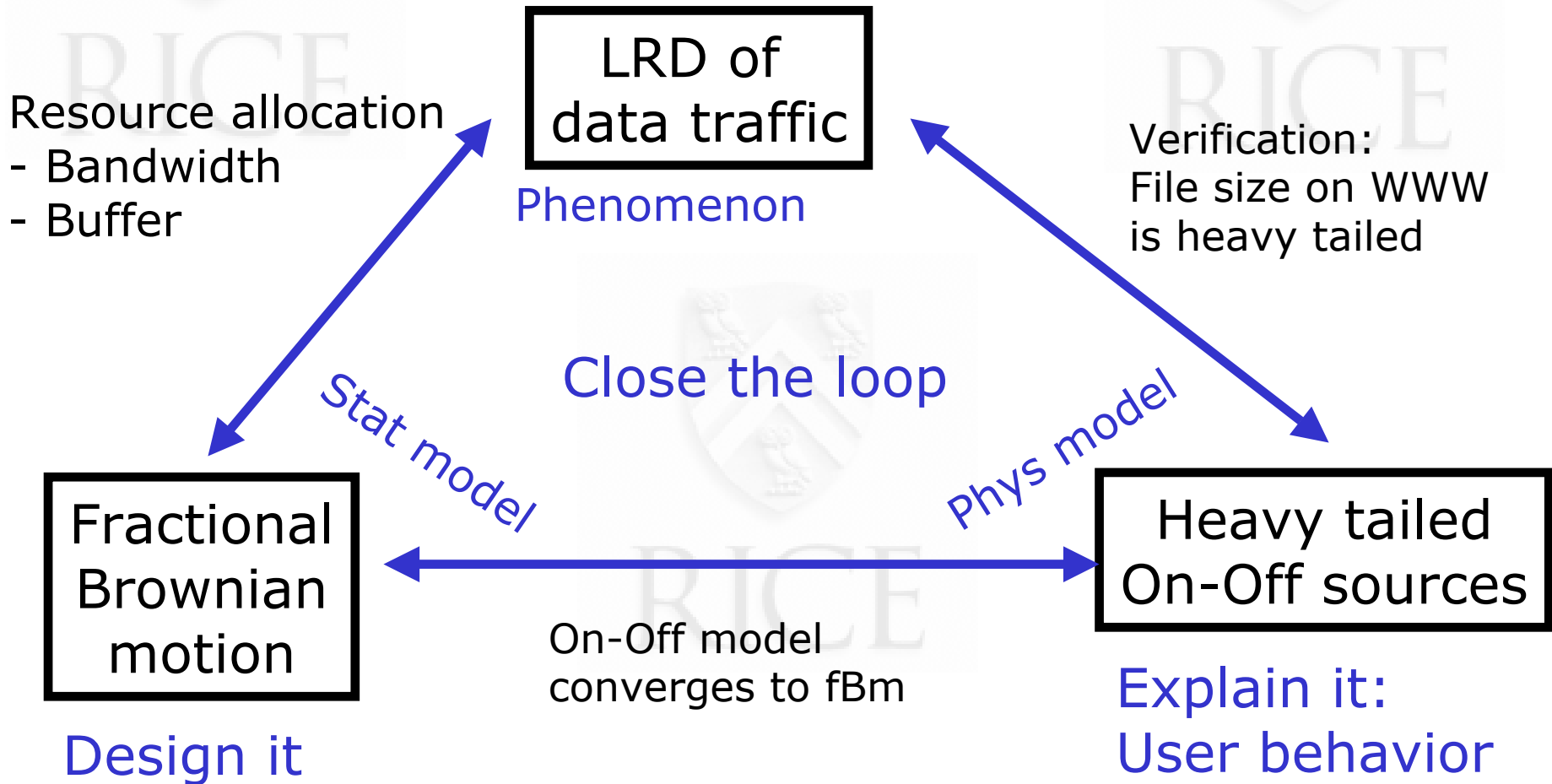
- Huge success

Voice: Heavy tails are rare

Data: Heavy tails are there



Understanding Large Scale





Small scale behavior

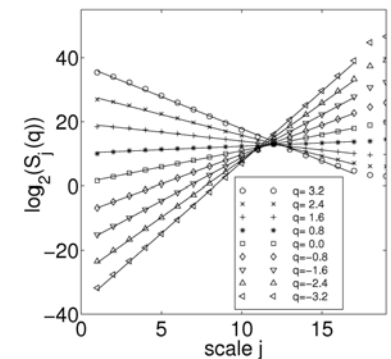
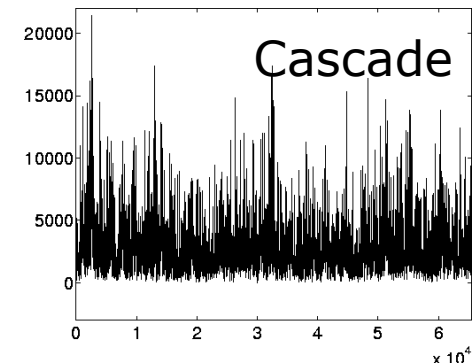
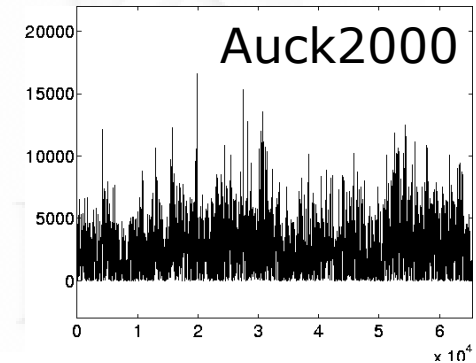


Below Round Trip Time

- fBm is **realistic** only at **large scales**
 - Explains how to **design**
- ...but is **unrealistic** at the small time scales relevant for
 - Control,
 - Performance (jitter, delay)
 - Service level agreements
 - Quality of Service (QoS)

Network Traffic is Multifractal

- Visually striking
- Scaling of impressive quality
(Levy Vehel & RR '96, Norros & Mannersalo '97, Willinger et al '98)
- Statistical models:
 - Binomial cascades
(Crouse & RR '98, Willinger et al '98)
 - On-off cascades (Baras)
 - Stationary Products
(Norros Mannersalo RR '99)
 - Products of Cylindrical Pulses
(Barral Mandelbrot '99)
 - IDC
(Bacry-Muzy '02, Chainais Abry RR)





Multiscale statistical modeling

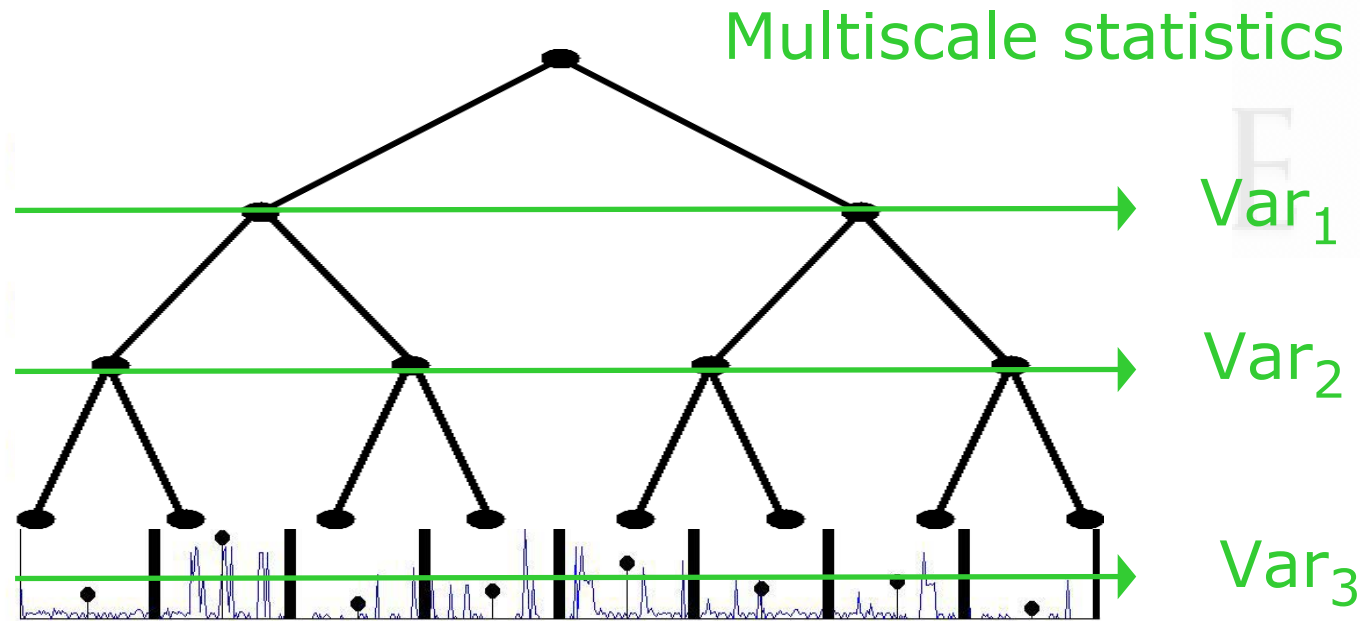
Why multiplicative?



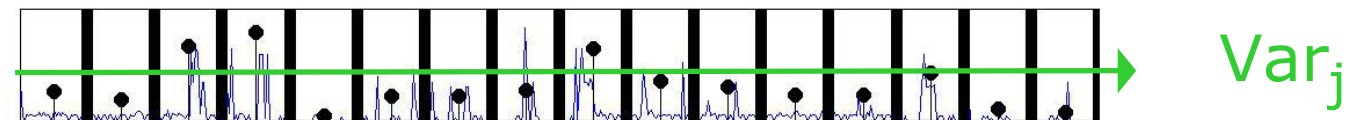
Multiscale Modeling

Time →

Scale ↓



Analysis: flow up the tree by adding



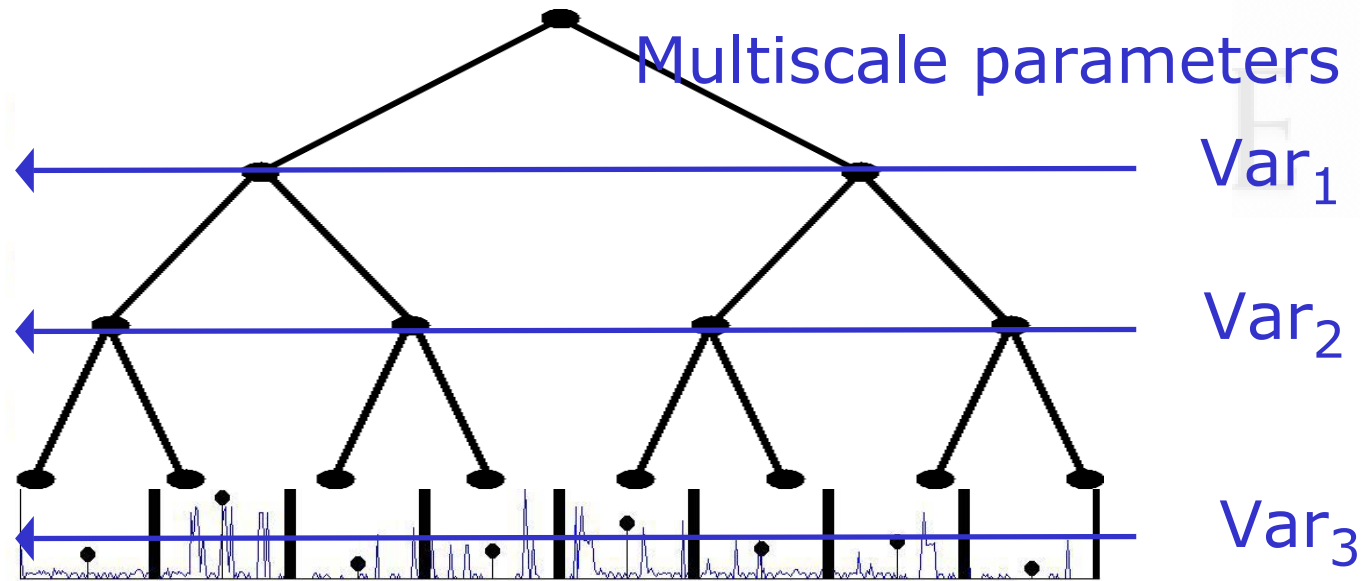
Start at bottom with trace itself

Multiscale Modeling

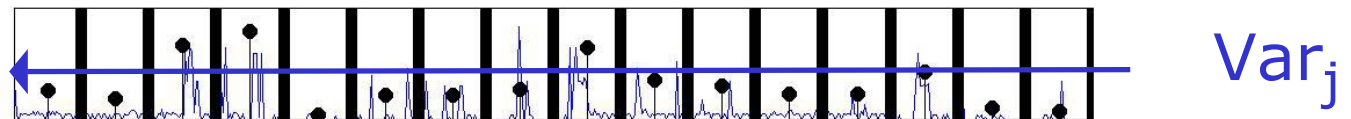
Time →

Start at top with total arrival

Scale ↓

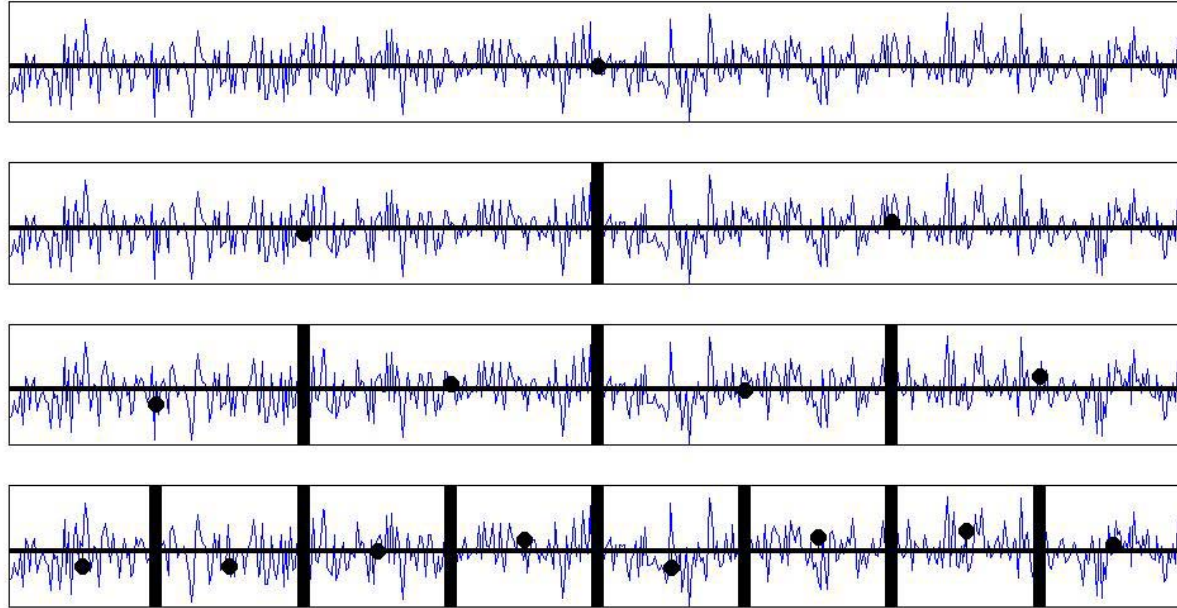


Synthesis: flow down via innovations



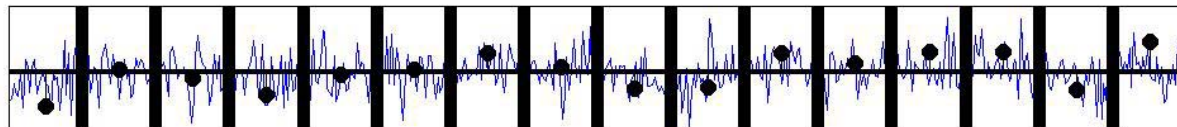
Signal: bottom nodes

Additive innovations: Linear Processes



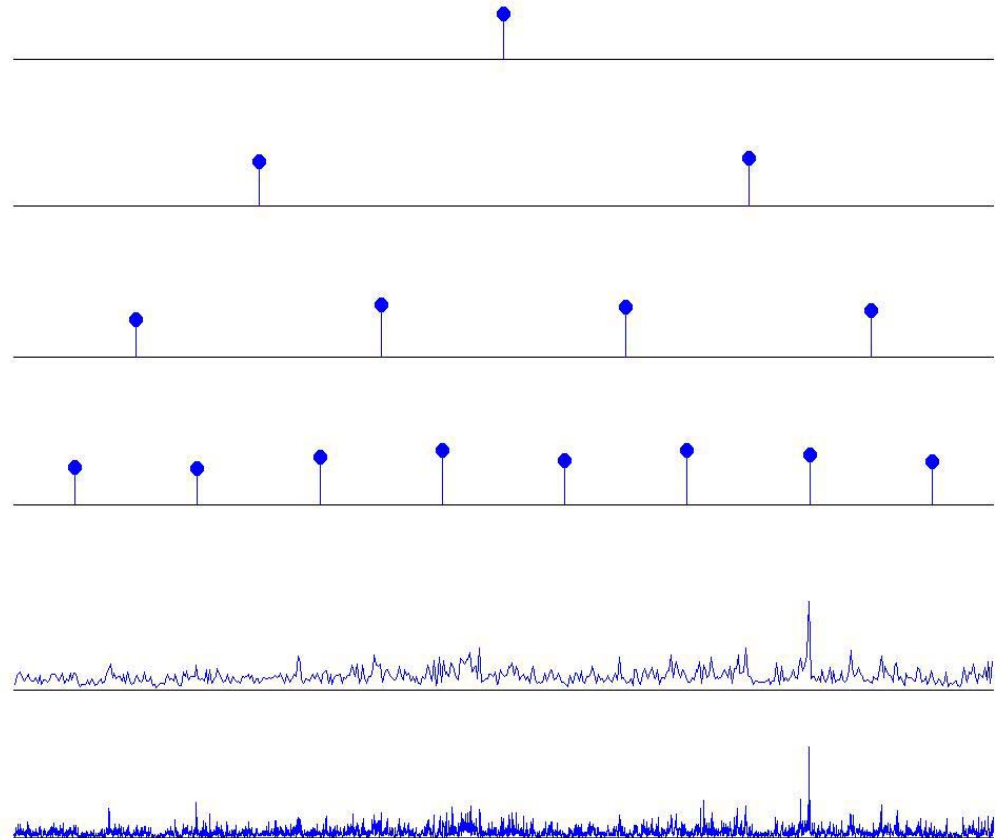
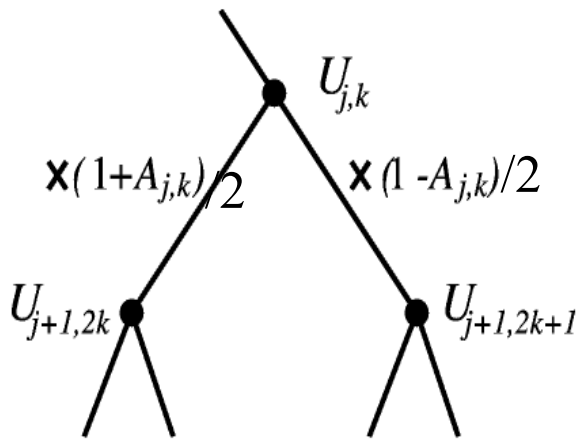
Match variances on all dyadic scales

CLT: asymptotically Gaussian



Additive Innovations $W_{jk} \sim \mathcal{N}(0, \sigma^2 2^{-j(2H-1)})$: Model for $B_H(t)$

Multiplicative Wavelet-Model: MWM



Multiplicative Innovations $(1 \pm A_{jk})/2 \sim \text{Beta}(1/2, \sigma_j)$: Variance-Match

Statistical relevance of cascades in networking

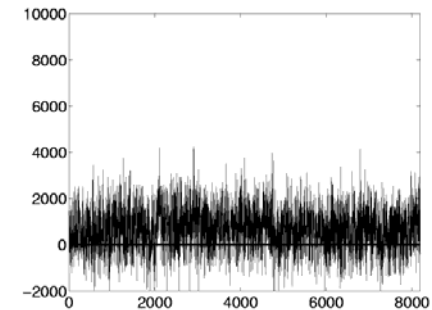
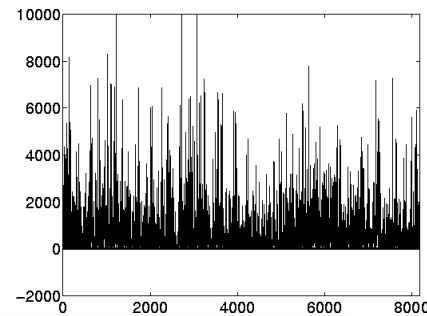
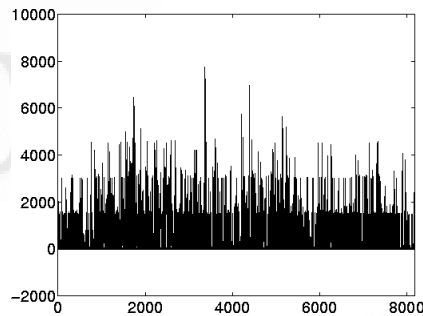
scale

Real trace (Auck00)

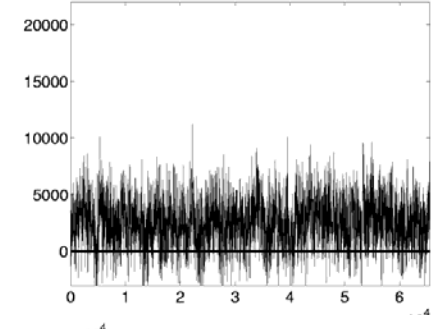
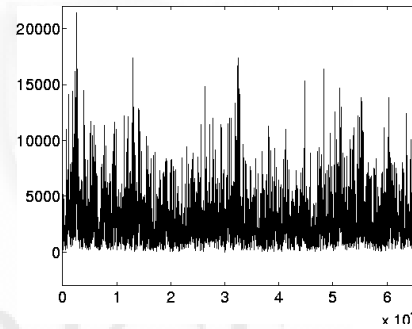
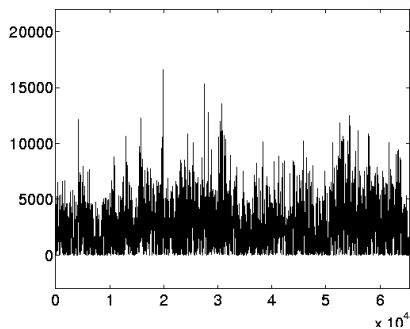
Cascade

Gaussian (fGn)

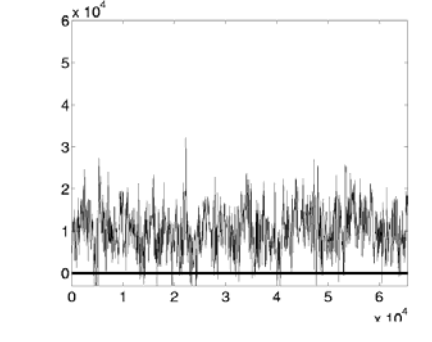
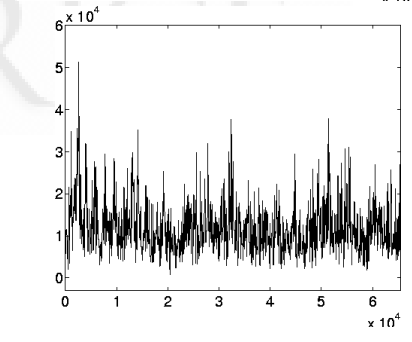
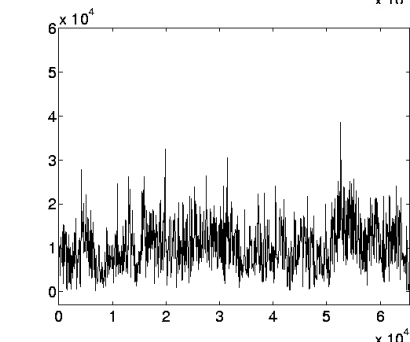
4ms



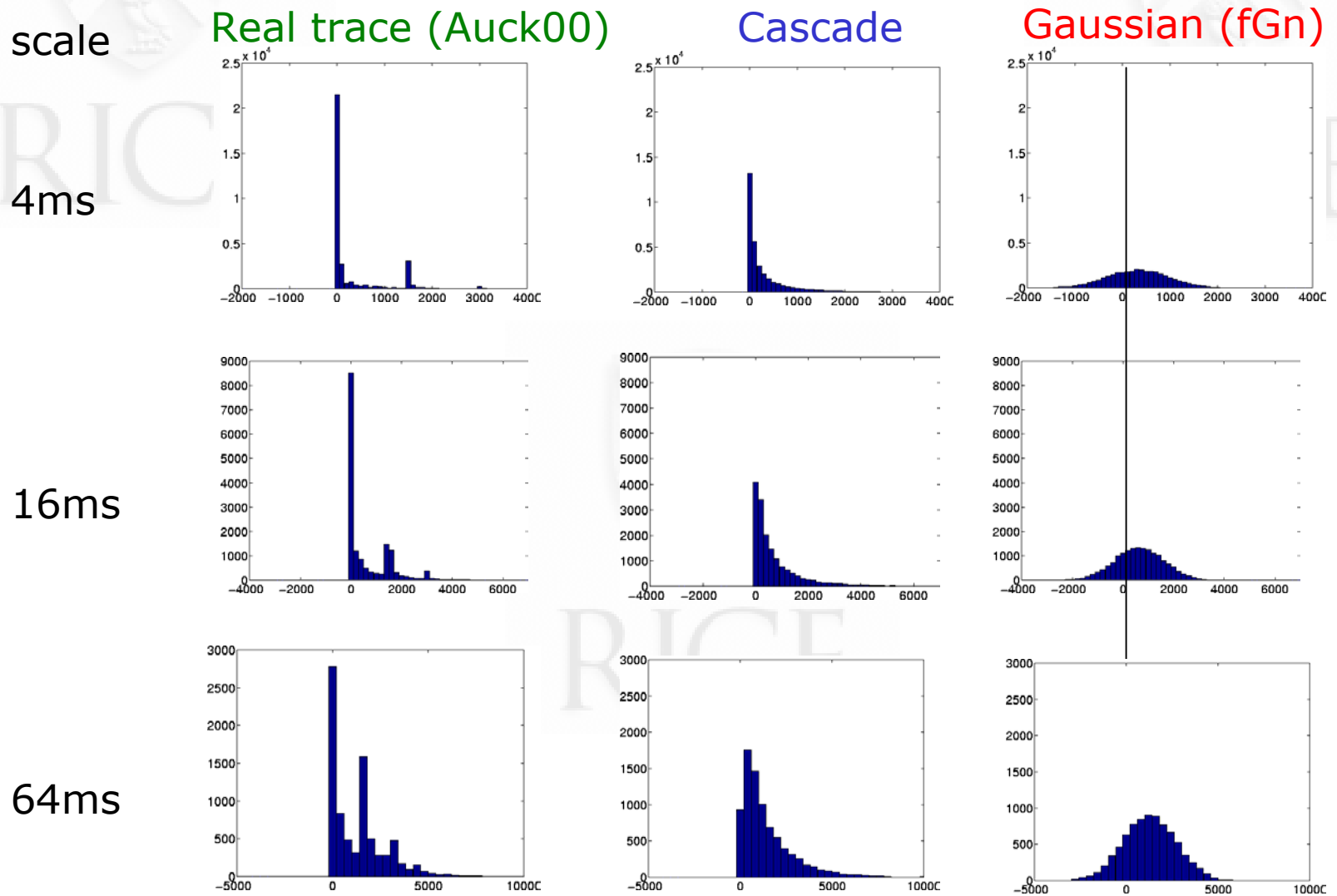
16ms



64ms



Statistical relevance of cascades in networking

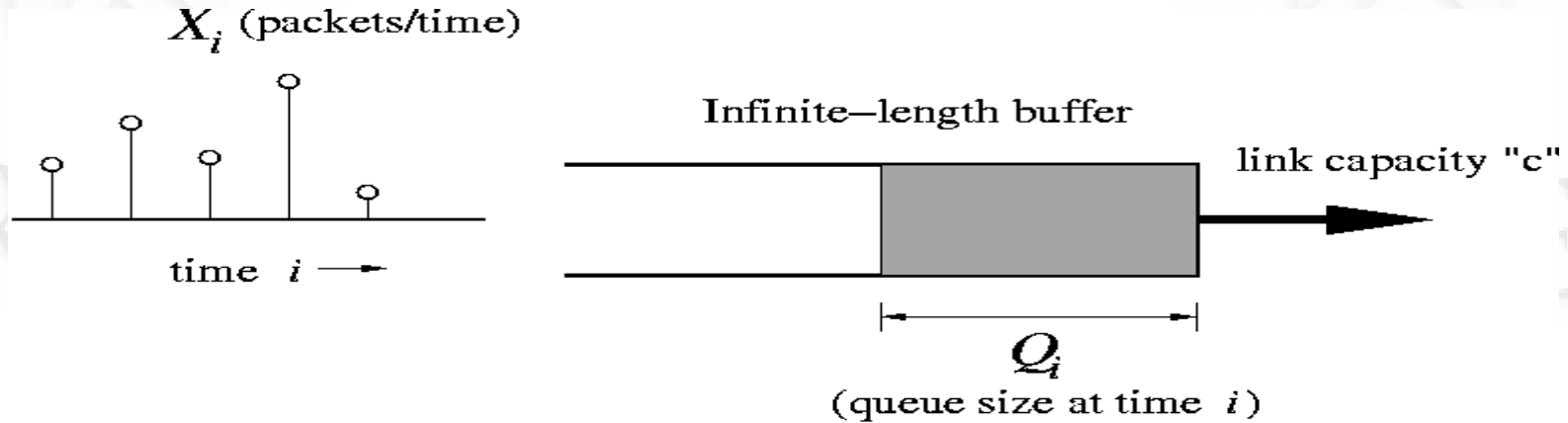




Own it:

**Queuing
Traffic inference**

Queuing 101



- Lindley's recursion

Q_i : queue size at time i

X_i : traffic arriving at time bin i

$$\begin{aligned}
 Q_0 &= \max(Q_{-1} + X_0 - c, 0) \\
 &= \max(\max(Q_{-2} + X_{-1} - c, 0) + X_0 - c, 0) \\
 &= \max(Q_{-2} + X_{-1} + X_0 - 2c, X_0 - c, 0)
 \end{aligned}$$



$$Q_0 = \max_n [X_1 + \dots + X_n - nc]$$

MultiScale Queuing approach

For tree models of traffic:

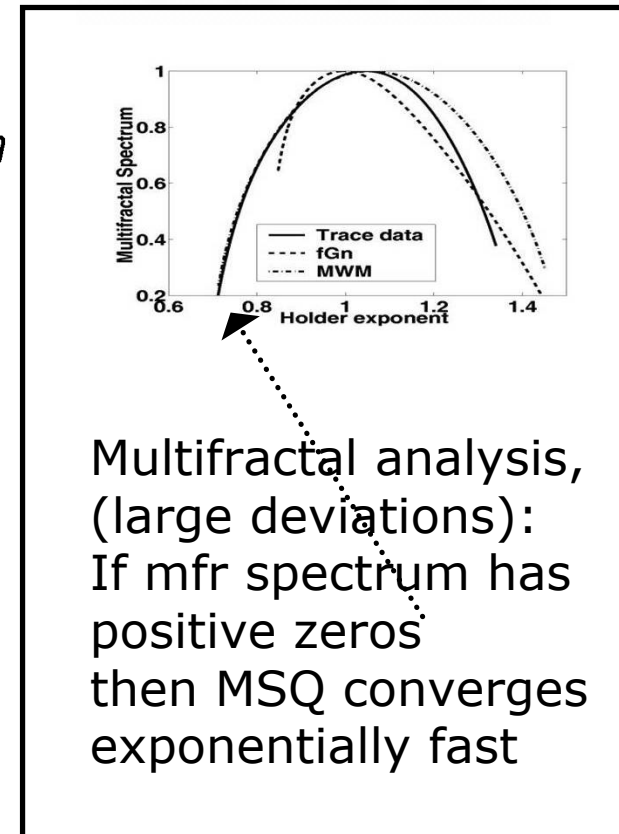
- Innovations W_i between **dyadic** times are **independent**
(pass to log for multiplicative model)

Lemma: $E_i := \{W_0 + \dots + W_{i-1} < b_i\}$,
 W_i are independent, otherwise arbitrary. Then

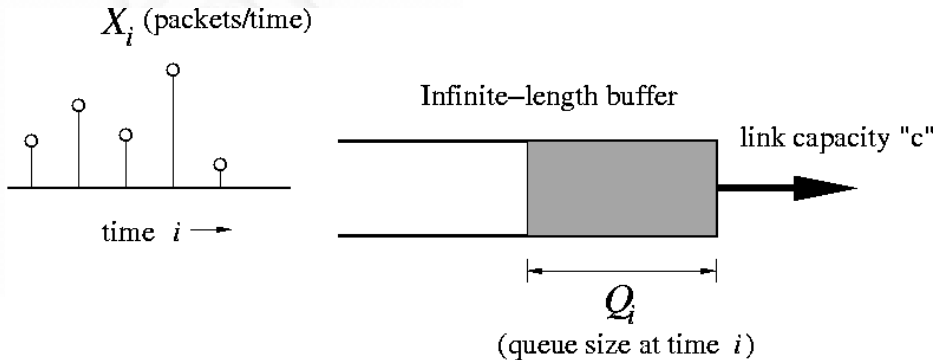
$$P[E_i | E_{i-1}, \dots, E_0] \geq P[E_i].$$

Dyadic queue tails

$$\begin{aligned} P[Q_D > b] &= 1 - P[Q_D < b] = 1 - P[\cap_{i=0}^n E_i] \\ &= 1 - P[E_0] \prod_{i=1}^n P[E_i | E_{i-1}, \dots, E_0] \\ &\leq 1 - \prod_{i=0}^n P[E_i] =: \text{MSQ}_n(b). \quad (1) \end{aligned}$$

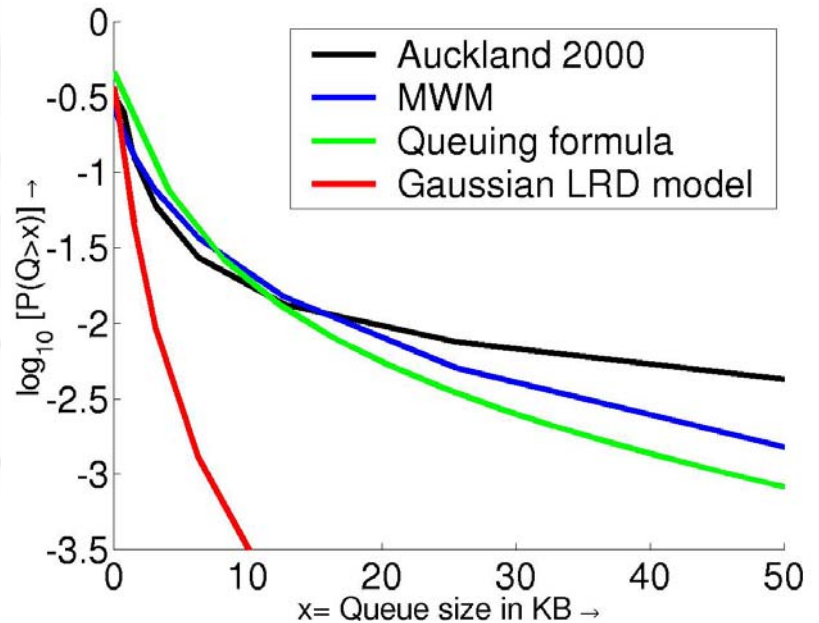


Queuing analysis



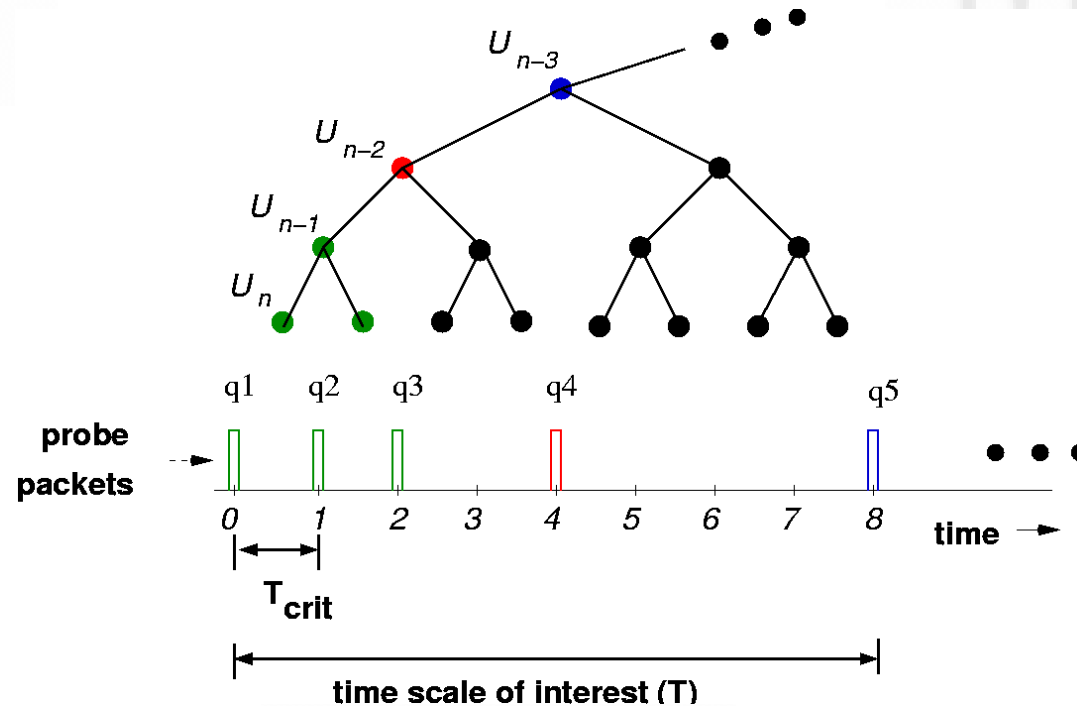
Q-tail: $P[Q > b]$

- Tree structure allows for **analytical queuing** formula
- **Multiplicative** model superior to **additive**
- Importance of multiscale tails (not only second order statistics)



Efficient Probing: *Packet Chirps*

- Tree inspires geometric **chirp probe**
- **MLE estimates** of cross-traffic at multiple scales



- Chirp is **practical** (ITC2000, Best paper PAM2003)
- Uniform is **optimal spacing** (Lehmann symposium 2004)



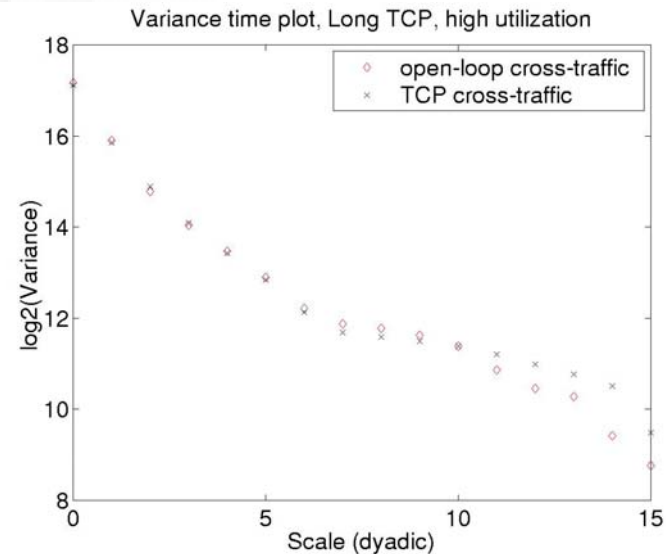
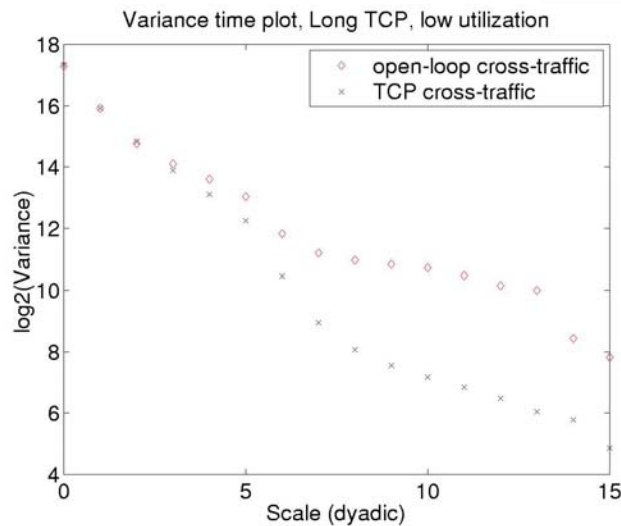
Statistical scaling at fine resolution



Beyond powerlaws

Real world data

- is stationary
- can deviate from powerlaws: traffic
- has no preference for dyadic scales



Beyond Self-similarity

- Self-similarity revisited:
 - $B(at) \stackrel{d}{=} C(a) B(t)$ B: process, C: scale function
 - $B(abt) \stackrel{d}{=} C(a)C(b) B(t)$
 - $C(a)C(b)=C(ab) \rightarrow C(a) = a^H$
 - $E[| B(a^n) |^q] = c(q) (a^{qH})^n$
 - linear in q (mono-fractal)

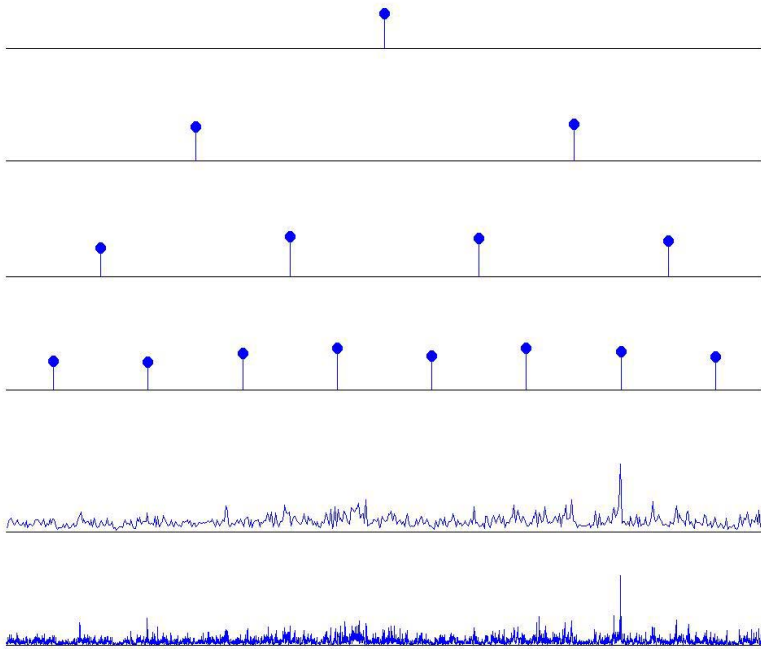
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 - linear in q (mono-fractal)
- More flexible rescaling “Ansatz”:
 - $C=C(a,t) ?$: non-stationary increments
 - C =independent r.v. for every re-scaling :
 - $X(a\dots at) = X(a^n t) = C_1(a)\dots C_n(a) X(t)$: **multiplicative**
 - $E[|X(a^n)|^q] = c(q) E[|C(a)|^q]^n$
 - non-linear in q (multifractal)

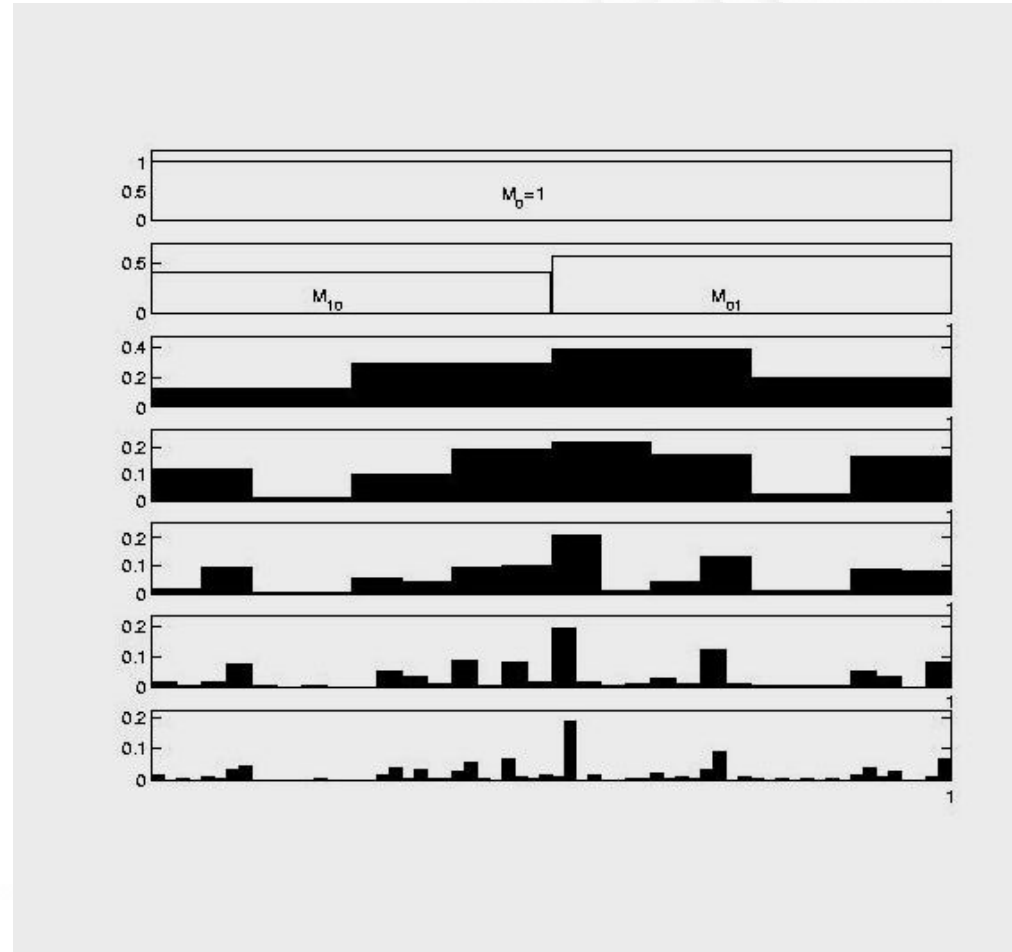
Beyond power-laws

- Self-similarity revisited:
 - $B(at) =^d C(a) B(t)$ B: process, C: scale function
 - $E[| B(a^n) |^q] = c(q) (a^{qH})^n$ (mono-fractal)
- More flexible rescaling “Ansatz”:
 - $X(a...at) = X(a^n t) = C_1(a)...C_n(a) X(t)$: multiplicative
 - $E[| X(a^n) |^q] = c(q) E[| C(a) |^q]^n$ (multifractal)
- Infinitely divisible scaling
 - $E[| X(a^n) |^q] = \exp(v(a^n) \xi(q))$
 - Multifractal for $v(x) = \log_a x$; mono-fractal for $\xi(q) = qH$
 - Interpretation:
 - $E[| X(a^n) |^q]$ is the Laplace trafo of $\log[X(a^n)]$
 - The distribution of $\log[X(a^n)]$ is an $v(a^n)$ -fold convolution
 - Thus $X(a^n)$ itself is in **distribution** the $v(a^n)$ -fold product of a unit multiplier W with $E[| W |^q] = \exp(\xi(q))$
 - First **model**: Product of Cylindrical Pulses (Mandelbrot Barral)

Binomial Measure



As a tree



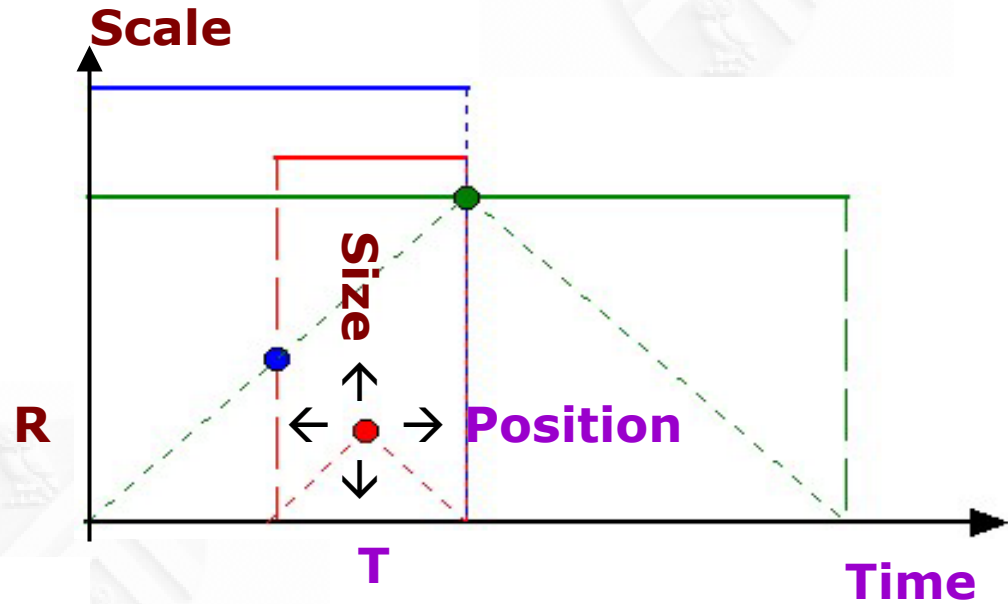
As a cascade of
multiplicative pulses

Geometry of Binomial Pulses

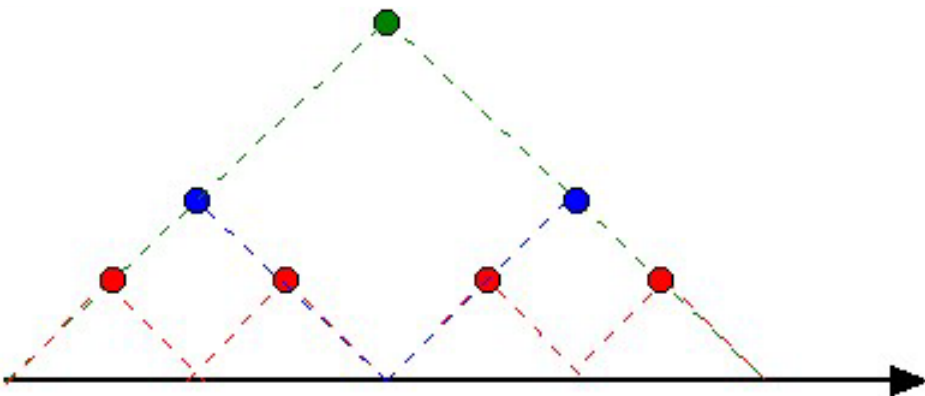
- Time-Scale plane: codes shape of pulses
 - **Position** (T =center)
 - **Size** (R =length)

Pulses:

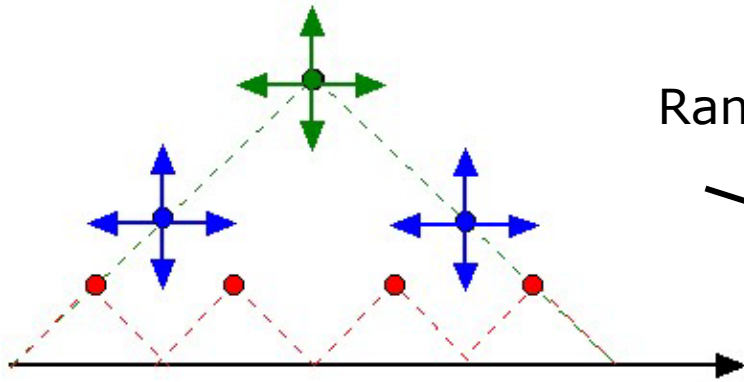
$$P_i(t) = \begin{cases} W_i & \text{if } |t-t_i| < r_i/2 \\ 1 & \text{else} \end{cases}$$



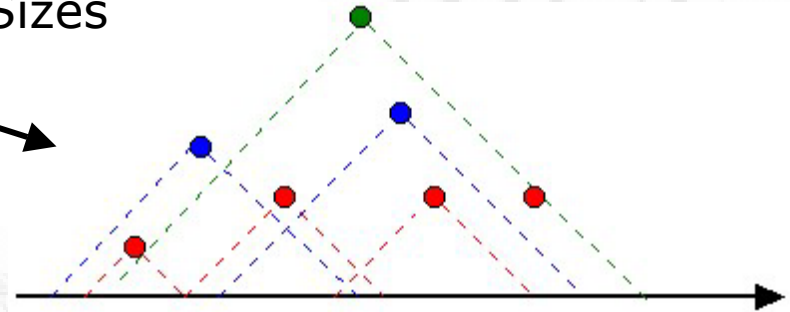
For Binomial:
Strict dyadic
geometry



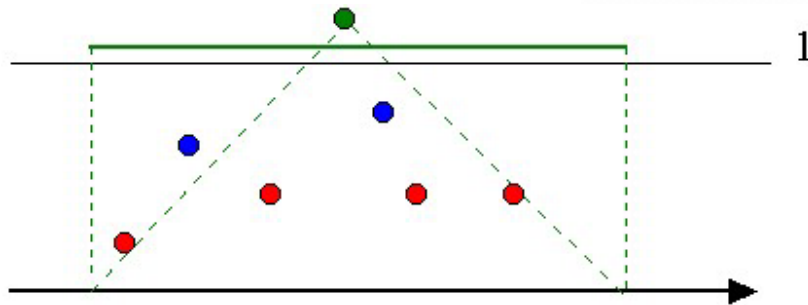
Stationary geometry



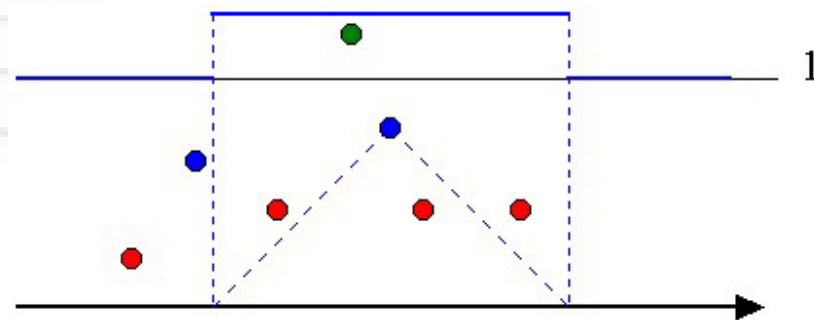
Randomize Positions
and Sizes



Large Scale Pulses



Medium Scale Pulses



Compound Poisson Cascade

Poisson points with control measure m

Cone of influence at t

$$C(t) = \{(t_i, r_i) : |t - t_i| < r_i/2\}$$

$$Q(t) = \prod \{W_i : |t - t_i| < r_i/2\}$$

Multifractal formalism

(self-similar case) Barral Mandelbrot

Multifractal Scaling

(non-powerlaws) (with Abry & Chainais)

$$T(q) = \exp[m(C(t,r))(1 - E[W^q])]$$

→ powerlaw only if $m(C(t,r)) = -\log(r)$

→ More general Infinitely Divisible Laws

